A Specification-based Test Generation Framework for RESTful Web Applications

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Overview

Before we start ...

Wep Applications

Webapps are everywhere!

Banking

Healthcare

eCommerce

ERP

eGovernance

Social
Networking

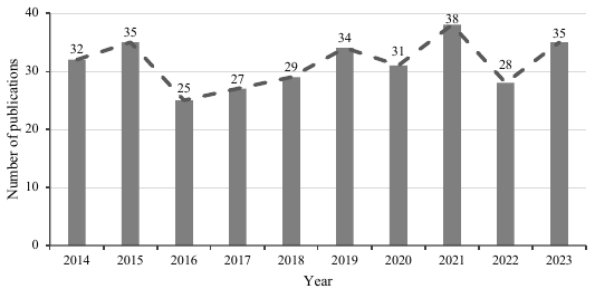
- Important
- Complex
- Distributed
- Dynamic

Desirable Properties

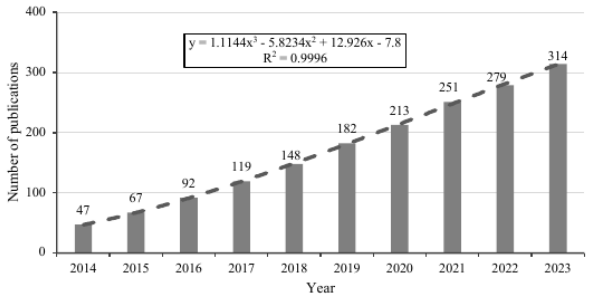
- **Functional properties:** Correct, complete, consistent
- **Non-function properties:** Performance (response time), scalable, available, reliable, secure, compliant, fair, inclusive, sustainable ...

Wep Applications

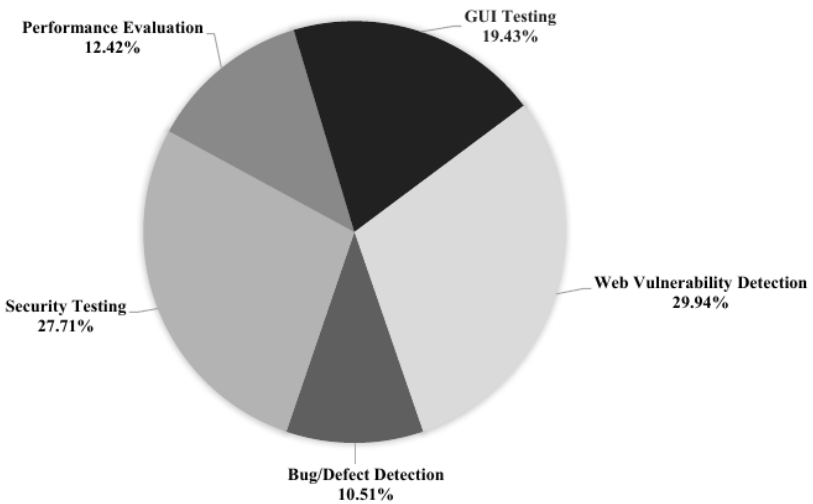
Testing of Webapps[1]



(a) Number of publications per year.



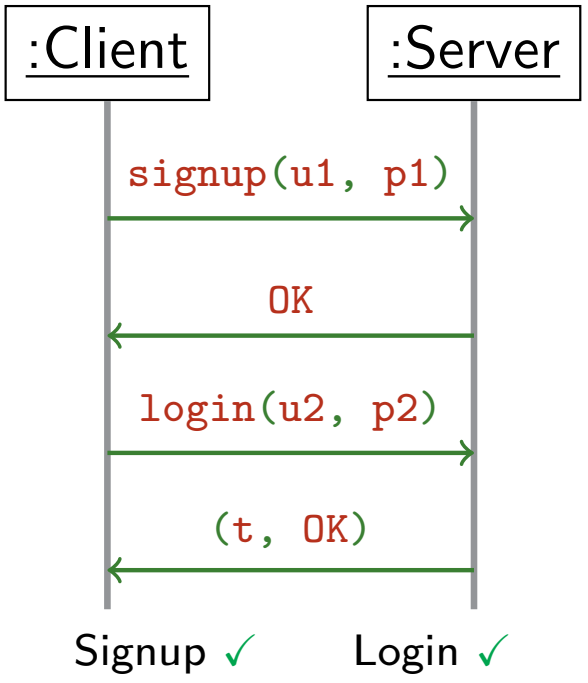
(b) Cumulative number of publications per year.

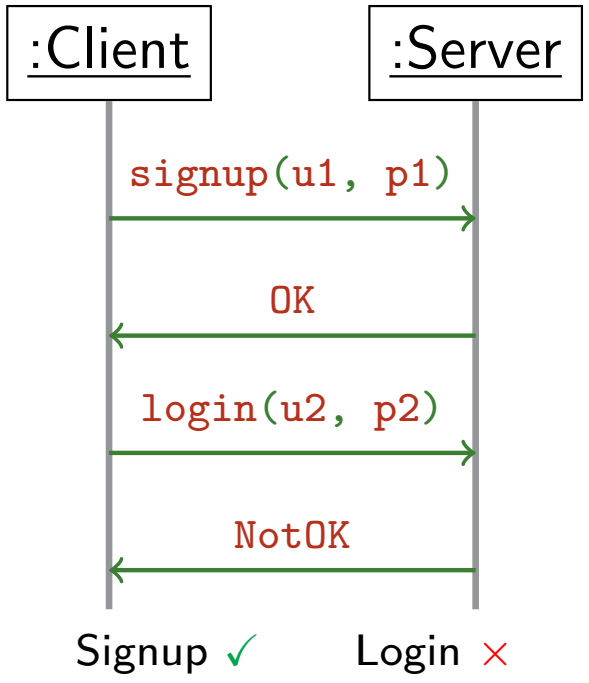


[1] A Survey on Web Application Testing: A Decade of Evolution - *Tao Li, Rubing Huang, Chenhui Cui, Dave Towey, Lei Ma, Yuan-Fang Li, Wen Xia*

Business Logic Testing

A Web Application





Business Logic Testing

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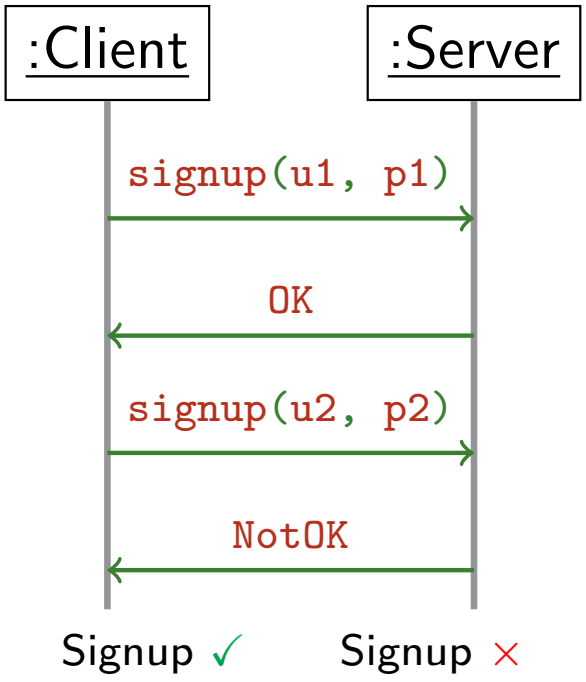

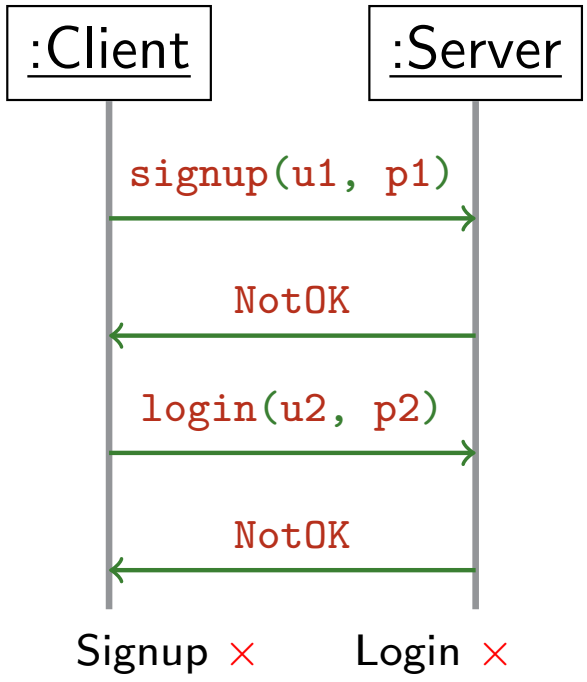

Business Logic Testing

A Web Application

Characteristics of business logic:

- Involves multiple client-server interactions
- Data dependency between interactions
- Dependency on continually evolving application state
- Too complex! Automation is essential.

Business Logic Testing

A Web Application

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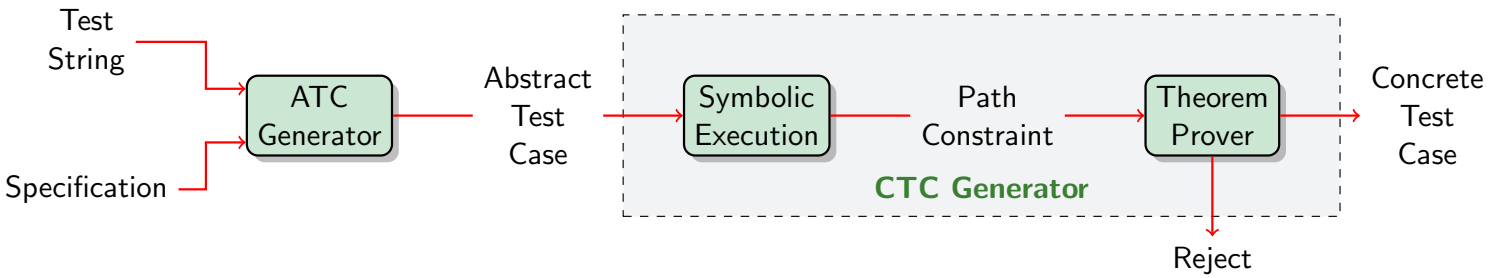
Our approach:

- Formally specify the system
- Automatically generate test cases from system + test specifications

Specification based Testing



Formal specification

Signup

- Before signup is called, the DB must not contain the username being currently used.
- A call to signup, if it succeeds, will return with an HTTP OK response code.
- After a successful signup, the DB will now have a record corresponding to the username that's been used for signing up.

Formal specification

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Formal specification

SignupOK	
Precondition	$u \notin dom(U)$
API	$\texttt{signup}(u, p) \rightarrow \mathit{HttpOK}$
Postcondition	$U' = U[u \mapsto p]$

Formal specification

Globals

$U : (string, string)map$
$T : (token, string)map$

Init

$U = \{\}$
$T = \{\}$

Functions

$signup : string \times string \rightarrow HTTPResponseCode$
$login : string \times string \rightarrow Token \times HTTPResponseCode$

SignupOK	
Precondition	$u \notin dom(U)$
API	$\texttt{signup}(u, p) \rightarrow HttpOK$
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LoginOK	
Precondition	$U[u] = p$
API	$\texttt{login}(u, p) \rightarrow (t, HttpOK)$
Postcondition	$T' = T[t \mapsto u]$

Formal specification

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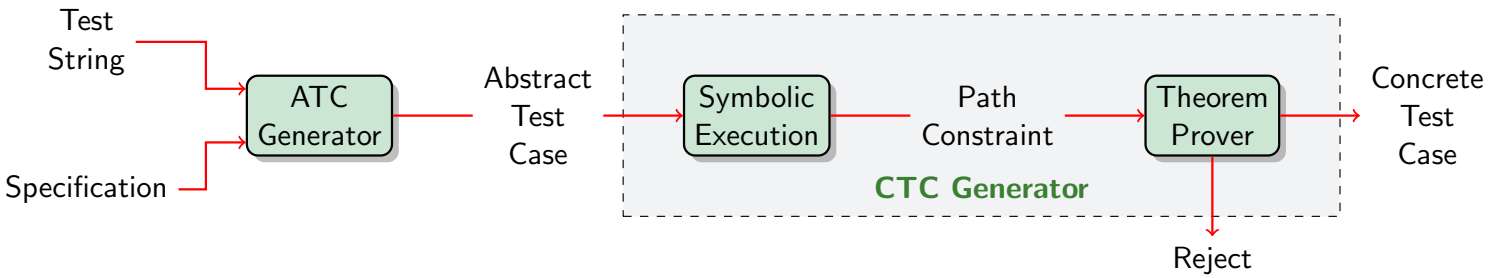
Specification based Testing



Abstract and Concrete Test Cases

Abstract and Concrete Test Cases

Abstract test case

```
let u1 := input<string>()
let p1 := input<String>()
let u2 := input<String>()
let p2 := input<String>()
let U = new Map<String, String>()
let T = new Map<Token, String>()

assume(u1 ∉ dom(U))
let r1 := signup(u1, p1)
assert(U[u1] = p1)

assume(u2 ∈ dom(U))
let (r2, t) := login(u2, p2)
assert(T[t] = u2)
```

Concrete test case

```
let u1 := "xyz"
let p1 := "abc"
let u2 := "xyz"
let p2 := "abc"
let U = new Map<String, String>()
let T = new Map<Token, String>()


let r1 := signup(u1, p1)
assert(U[u1] = p1)


let (r2, t) := login(u2, p2)
assert(T[t] = u2)
```

The GETATC Algorithm

Test string: `f f`

Formal Specification:

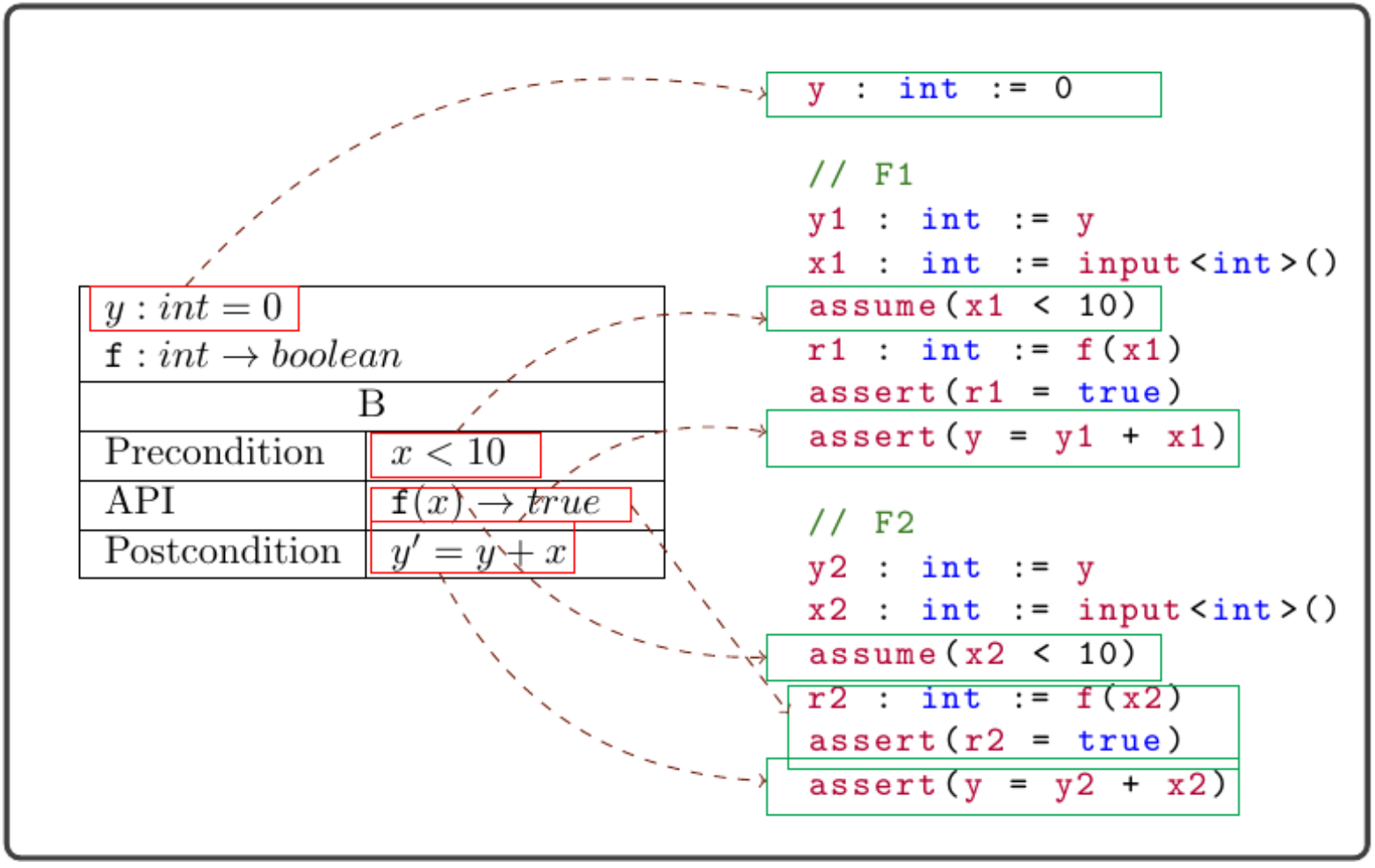
$y : int = 0$ $\texttt{f} : int \rightarrow boolean$	
Precondition	$x < 10$
API	$\texttt{f}(x) \rightarrow true$
Postcondition	$y' = y + x$

```
y : int := 0

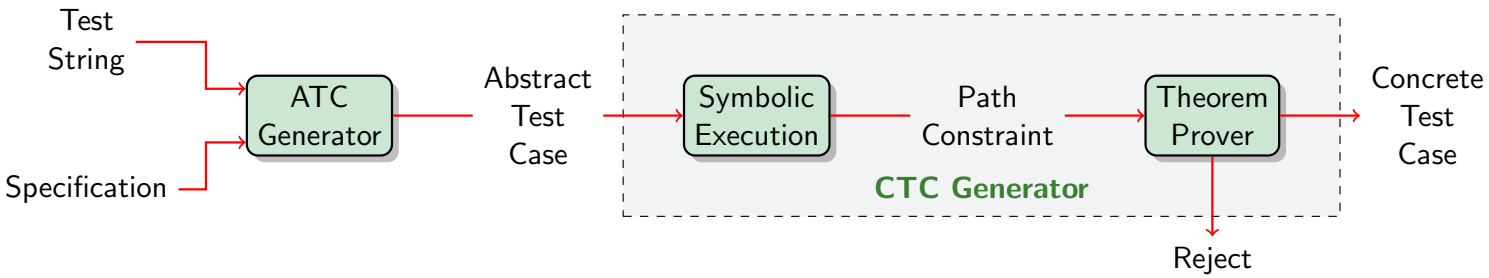
// F1
y1 : int := y
x1 : int := input<int>()
assume(x1 < 10)
r1 : int := f(x1)
assert(r1 = true)
assert(y = y1 + x1)

// F2
y2 : int := y
x2 : int := input<int>()
assume(x2 < 10)
r2 : int := f(x2)
assert(r2 = true)
assert(y = y2 + x2)
```

The genATC Algorithm

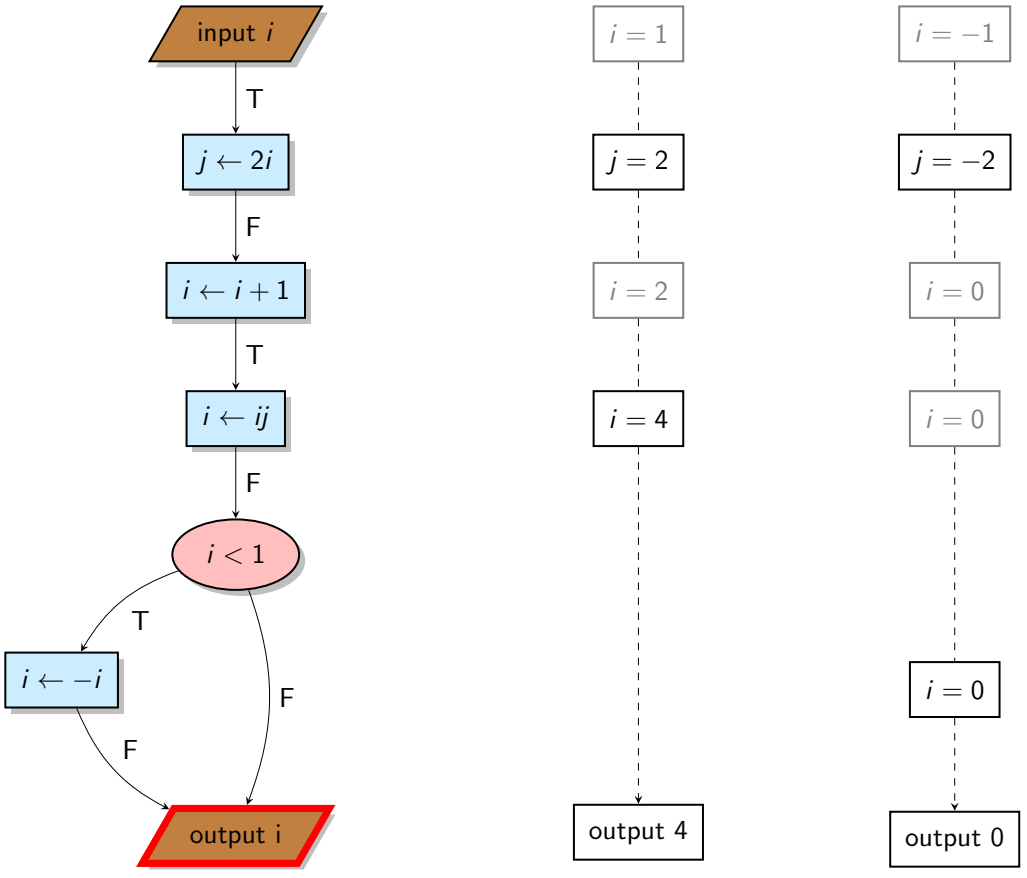


Specification based Testing



Symbolic Execution

Concrete Execution – Example



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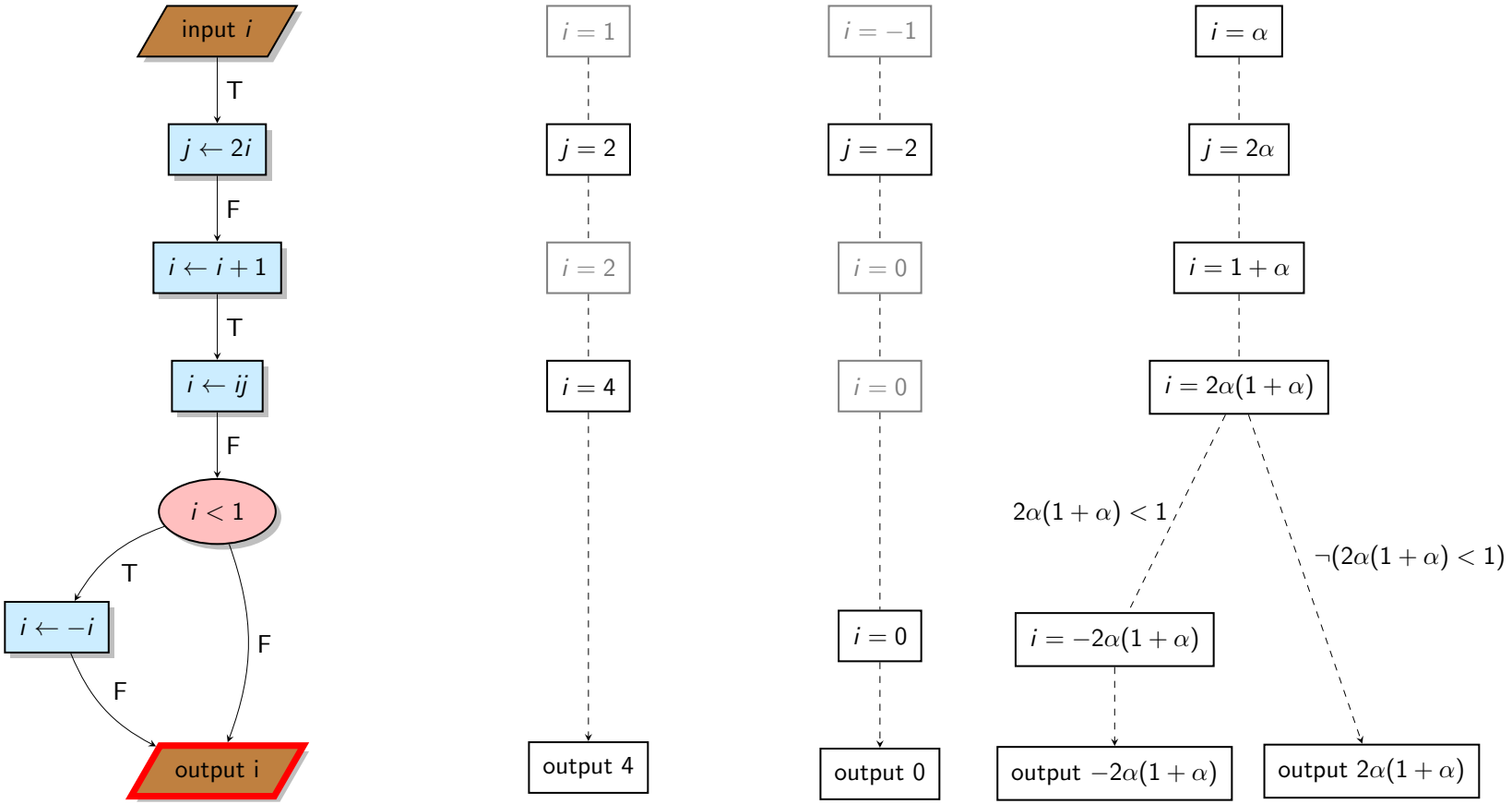

Generating Concrete Test Case

```
function genCTC(t, L, σ)
    if ¬isAbstract(t) then
        return t
    else
        t' ← rewriteATC(t, L)
        L' ← symex(t', σ)
        return getCTC(t', L', σ)
```

- isAbstract(t): Returns true if the testcase t has atleast one input command.
- σ: Type environment

Rewriting Abstract Test Case

Considering the abstract test case as a sequence of statements:
$t = [s_1; s_2; ...; s_n]$ and $L = [v_1; v_2; ...; v_m]$

function REWRITEATC(t, L)
  if $|t| = 0 \wedge |L| \neq 0$ **then** raise Error
  match s_1 with
  | case $Input(x) \Rightarrow$
    $s_1' \leftarrow Assign(x, v_1)$
    return s_1' :: REWRITEATC($[s_2; ...; s_n]$ $[v_2; ...; v_m]$)
  | _ $\Rightarrow$ return s_1 :: REWRITEATC($[s_2; ...; s_n]$ $[v_1; ...; v_m]$)

Here :: is list construction operator. For example: $1 :: [2; 3] = [1; 2; 3]$

Rewriting Abstract Test Case

Example

t	L	t'
`x := input()`	1	`x := 1`
`y := input()`	2	`x := 2`
`z := input()`		`z := input()`
...		

Symbolic Execution

Considering the abstract test case as a sequence of statements:
$t = [s_1; s_2; ...; s_n]$

function SYMEX($[s_1, s_2, ..., s_n]$, σ)
    $C \leftarrow []$
    for $i = 1$ to n **do**
        if isReady(s_i) **then**
            symexInstr(s_i, σ, C)
        else
            break
    $pc \leftarrow$ computePathConstraint(C)
    return solve(pc)

Symbolic Execution

What do we mean by Ready?

$f(x, y),\ \sigma = [x \mapsto 1, y \mapsto 2]$	✓
$f(x, y),\ \sigma = [x \mapsto Add(1, X_1), y \mapsto 2]$	×

Symbolic Execution

isReady Function

function $\textsc{isReady}(s, \sigma)$
    match s with
    | case $\mathit{Assign}(x, e) \Rightarrow$ return $\textsc{isReady}(e, \sigma)$
    | ...

function $\textsc{isReady}(e, \sigma)$
    match e with
    | case $\mathit{Var}(x) \Rightarrow$ return $\neg\ \textsc{isSymbolic}(\sigma[x])$
    | case $\mathit{Num}(n) \Rightarrow$ return true
    | case $\mathit{FunCall}(f, [a_1; ...; a_n])$ where $\textsc{isAPI}(f) \Rightarrow$ return $\bigwedge_{i=1}^{n} \textsc{isReady}(a_i)$
    | case $\mathit{FunCall}(f, [a_1; ...; a_n])$ where $\neg\textsc{isAPI}(f) \Rightarrow$ return true
    | ...

Symbolic Execution

isSymbolic Function

```
function isSymbolic(e)
    match e with
    |   case Var(_) ⇒ return true
    |   case Num(n) ⇒ return false
    |   case Add(e1, e2) ⇒ return isSymbolic(e1) ∨ isSymbolic(e2)
    |   ...
```

Symbolic evaluation

Instruction

function SYMEXINSTR(s, σ, C)
  match s with
  | case *Assign*$(x, e) \Rightarrow \sigma \leftarrow \sigma[x \mapsto$SYMEVAL$(e, \sigma)]$
  | case *Assume*$(c) \Rightarrow C \leftarrow C \frown$ SYMEVAL(c, σ)
  | ...

Here, $\frown$ stands for addition of an element to a list:

$[1; 2; 3] \frown 4 = [1; 2; 3; 4]$

Symbolic Execution

Computing the path constraint

Expression

function SYMEVAL(e, σ)
    match e with
    | case $Var(x) \Rightarrow$ return $\sigma[x]$
    | case $Num(n) \Rightarrow$ return $\sigma[x]$
    | case $Add(e_1, e_2) \Rightarrow$ return $Add(\text{SYMEVAL}(e_1, \sigma[x]), \text{SYMEVAL}(e_2, \sigma[x]))$
    | ...

Computing Path Constraint

function COMPUTEPATHCONSTRAINT($[c_1, c_2, ..., c_n]$)
    return $c_1 \wedge c_2 \wedge ... \wedge c_n$

The TestGen Algorithm

Symbolic Execution – Run 1

```
y : int := 0

// F1
y1 : int := y
x1 : int := input<int>()
assume(x1 < 10)
r1 : int := f(x1)
assert(r1 = true)
assert(y = y1 + x1)

// F2
y2 : int := y
x2 : int := input<int>()
assume(x2 < 10)
r2 : int := f(x2)
assert(r2 = true)
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```

1. $y = 1$
2. $y1 = 1$
3. $x1 = X_1$
4. assume$(X_1 < 10)$

The TestGen Algorithm

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f needs a concrete value of $x1$!

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```

1. $y = 1$
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f needs a concrete value of $x1$!

Solution:

1. Compute path constraint so far: $X_1 < 10$.
2. Give to SMT solver to solve. SAT. Example value: $X_1 = 5$.

The TestGen Algorithm

Abstract Test Case (version 2)

```
y : int := 0
y1 : int := 0
x1 : int := 5
r1 : int := f(x1)
...
...
```

The TestGen Algorithm

Symbolic Execution – Run 2

1. $y = 1$
2. $y1 = 1$
3. $x1 = 5$
4. $f(x1) \rightarrow v_1$
5. $r1 = v_1$
6. assert$(r1 = \mathit{true})$
7. assert$(y = y1 + 5)$
8. $y2 = y$
9. $x2 = X_2$
10. assume$(X_2 < 10)$

The TestGen Algorithm

Symbolic Execution – Run 2

1. $y = 1$
2. $y1 = 1$
3. $x1 = 5$
4. $f(x1) \rightarrow v_1$
5. $r1 = v_1$
6. assert$(r1 = \mathit{true})$
7. assert$(y = y1 + 5)$
8. $y2 = y$
9. $x2 = X_2$
10. assume$(X_2 < 10)$

f needs a concrete value of $x2$!

The TestGen Algorithm

Symbolic Execution – Run 2

1. $y = 1$
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8. $y2 = y$
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10. assume$(X_2 < 10)$

f needs a concrete value of $x2$!

Solution:

1. Compute path constraint so far: $X_2 < 10$.
2. Give to SMT solver to solve. SAT. Example value: $X_2 = 2$.

The TestGen Algorithm

Abstract Test Case (version 3)

```
y : int := 0
y1 : int := 0
x1 : int := 5
r1 : int := f(x1)
assert(r1 = true)
assert(y = y1 + x1)
y2 : int := y
x2 : int := 2
r2 : int := f(x2)
assert(r2 = true)
assert(y = y2 + x2)
```

Implementation and Experiments

The Team

Research Scholars

Aira Jain

Pranita Ganguly

The Team

Research Scholars

Aira Jain

Pranita Ganguly

Implementation and Experiments

1. Prototype: Implemented using C++
2. SMT solver: Z3
3. Case studies: Student and Institute projects done by people outside the team
4. Test strings generated manually
5. Result: Successfully generated 100s of integration test cases
6. Fault injection
7. Our tests detect these faults while state-of-the-art REST testing tool can't.

Mock Generation using Contracts

Mocking – Problem

- f is the function under test (client side).
- f calls another function g (server side).
- g is a potentially side-effectful function that requires some prerequisite code to run before it becomes valid to call g.
- Disadvantages:
  1. Complex to prepare unit test for f (because it has to prepare the server for the test).
  2. Makes tests run slower.
  3. Sometimes, may not be feasible due to unavailability of server.

Mock Generation using Contracts

Motivation - Mocking

- Mocking allows us to simplify the problem of writing unit tests for f.
- Instead of using g, we use g', a mock of g.
- g' has the same function signature as g.
- g' is much simpler than g.
- g' returns a correct value that allows f's test to proceed even without our having to run the prerequisite code for g.

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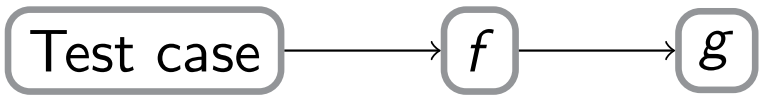
- g' is much simpler than g.
- g' returns a correct value that allows f's test to proceed even without our having to run the prerequisite code for g.
- g' runs on the client.

Testing without mocking

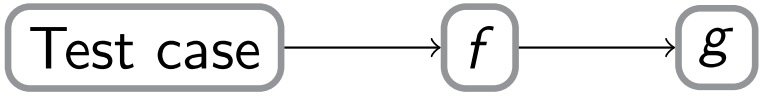


Test case

Prerequisite code for g

call to g

Testing without mocking



Test case

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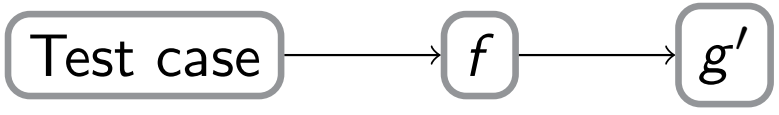
call to g

Testing with mocking



Test case

No Prerequisite code for g

call to g'

Mock Generation using Contracts

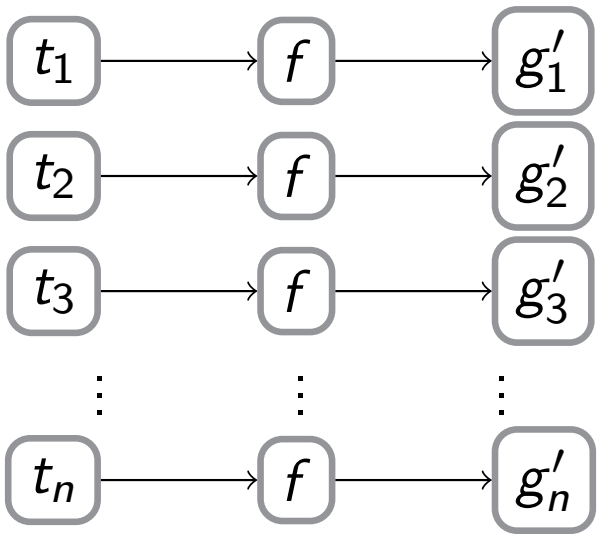
Mocking - Challenges

1. Ensuring valid precondition
2. Ensuring valid postcondition

Mock Generation using Contracts

Mocking - Challenges

1. Ensuring valid precondition
2. Ensuring valid postcondition
3.



Mock Generation using Contracts

Mocking - Challenges

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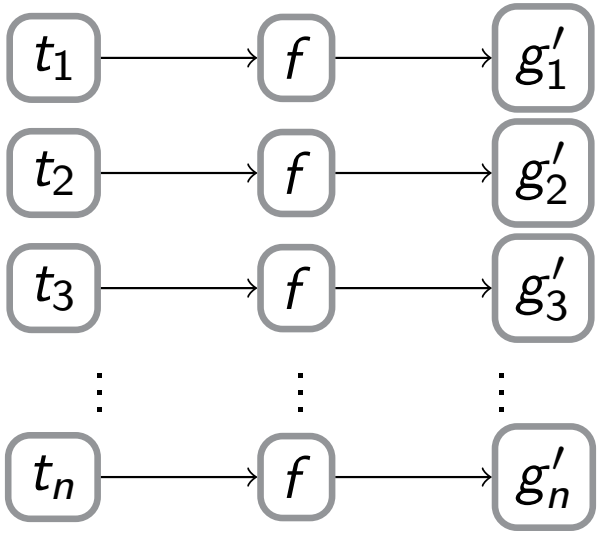

A DIFFERENT g' FOR EACH TEST CASE!

Mock Generation using Contracts

Mocking - Challenges

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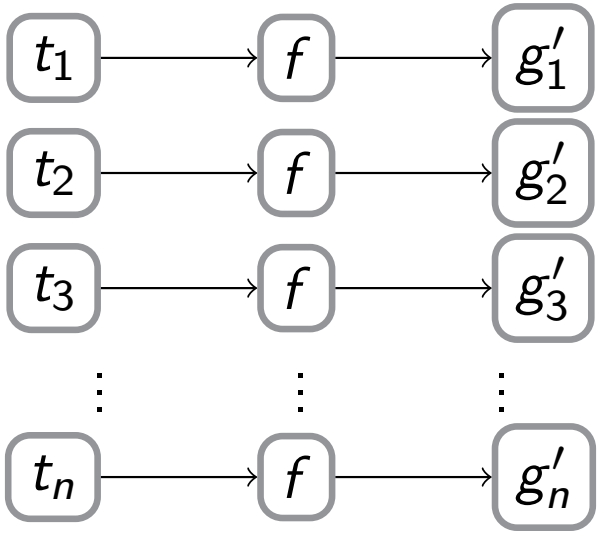
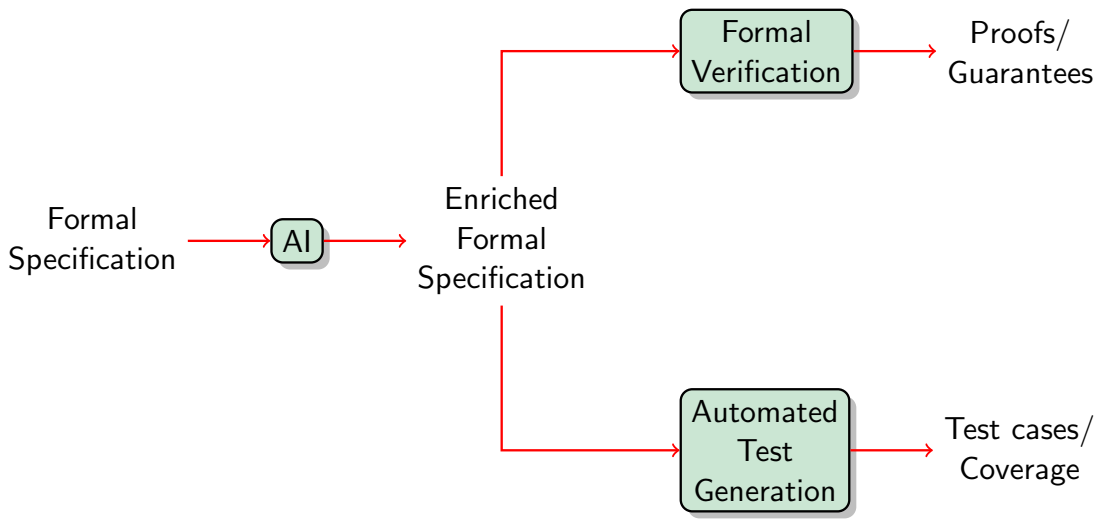


A DIFFERENT g' FOR EACH TEST CASE! – Automation necessary.

Specification Generation using LLMs



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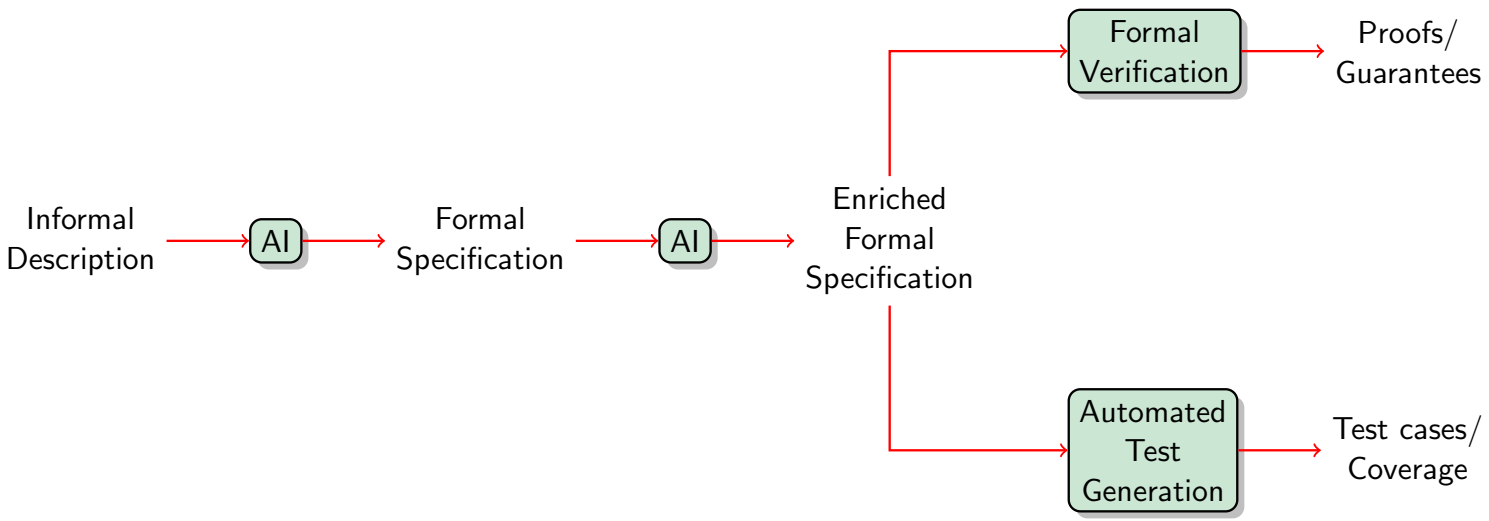
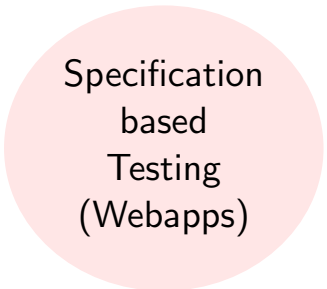


The TestGen Algorithm

A Roadmap



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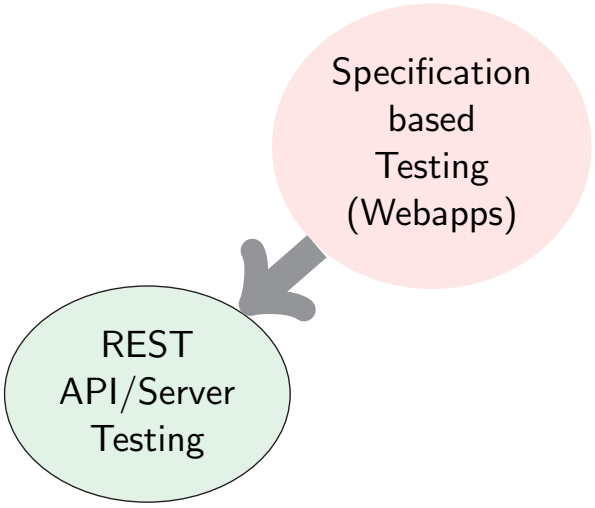

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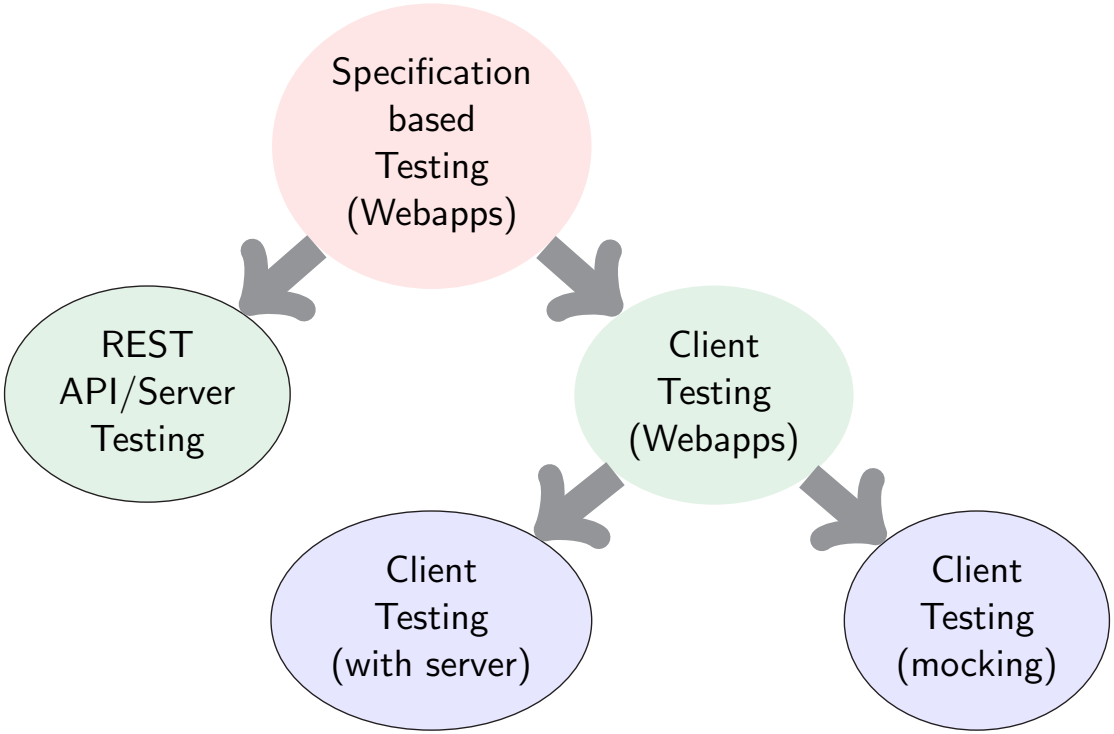
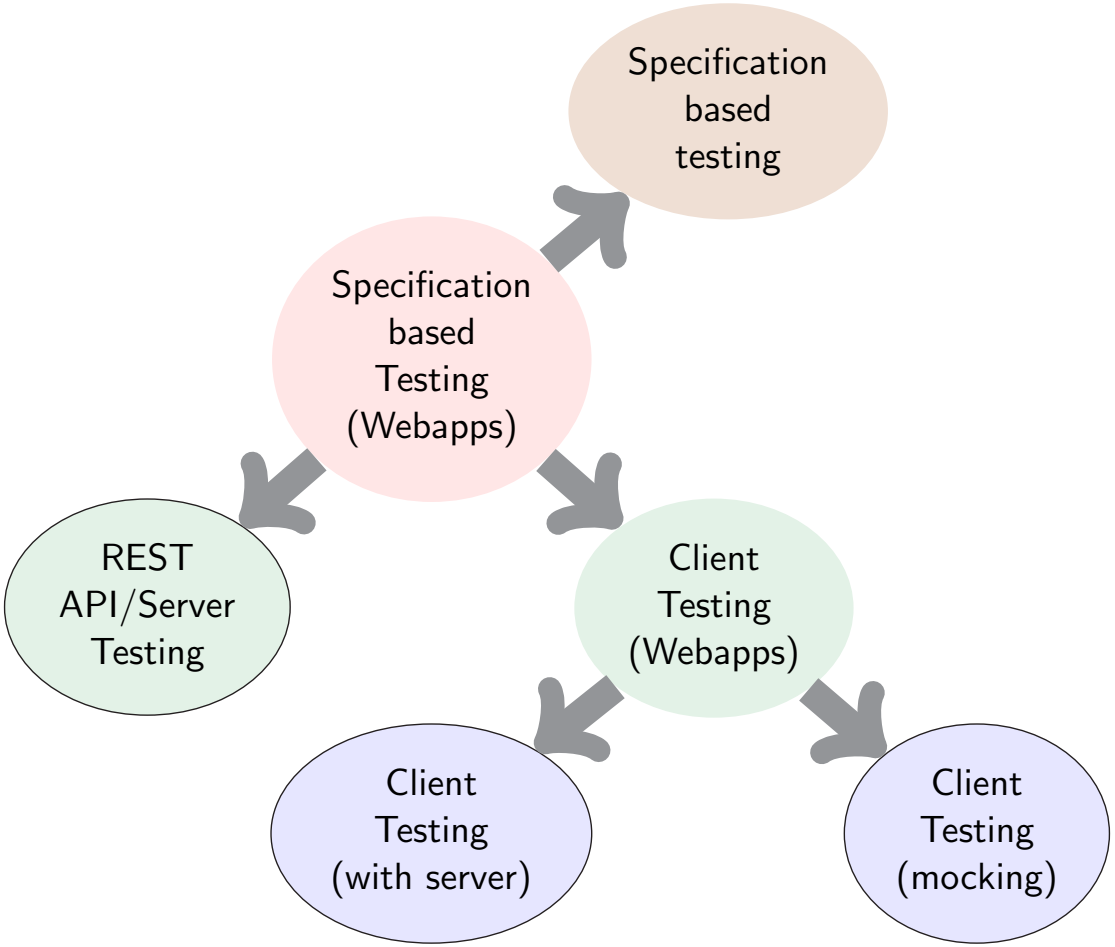


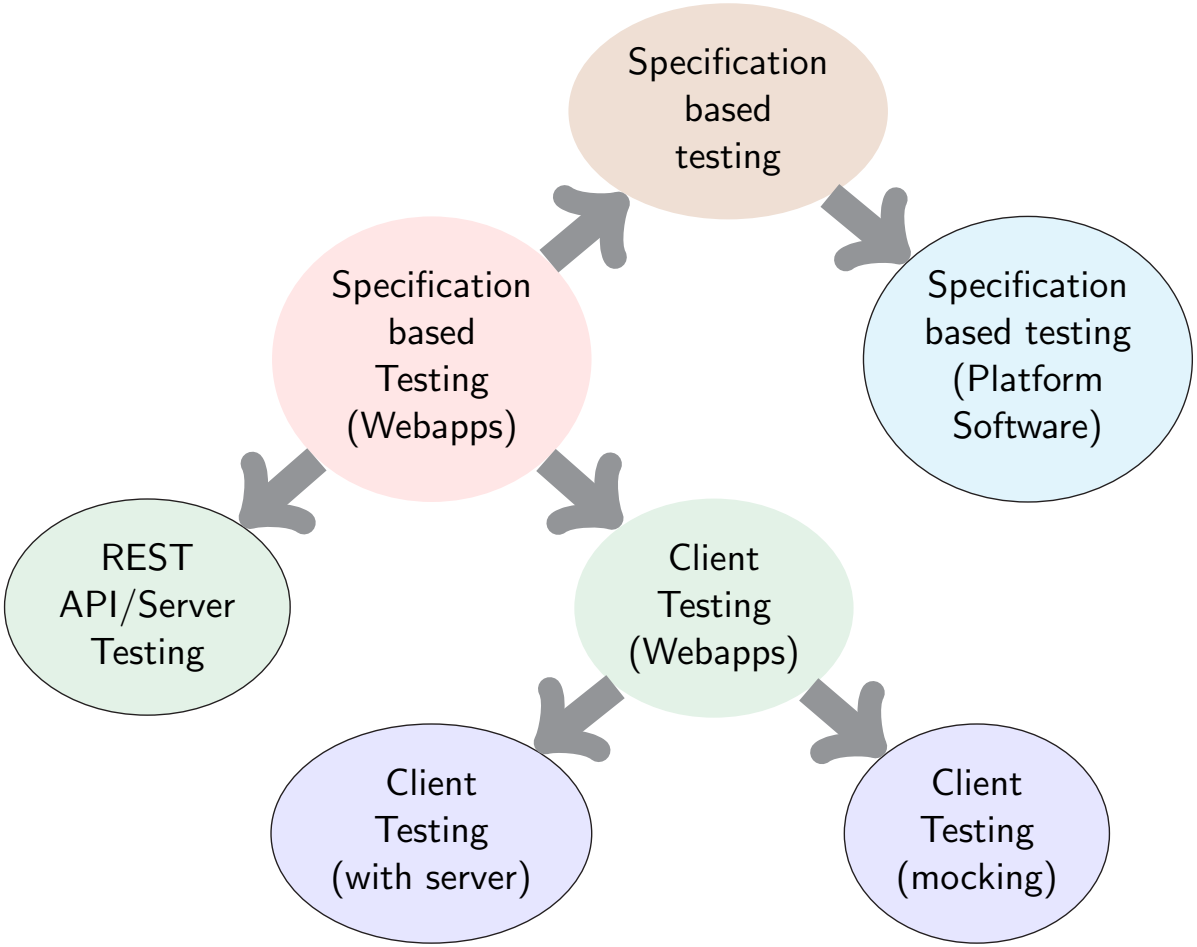
The TestGen Algorithm

A Roadmap



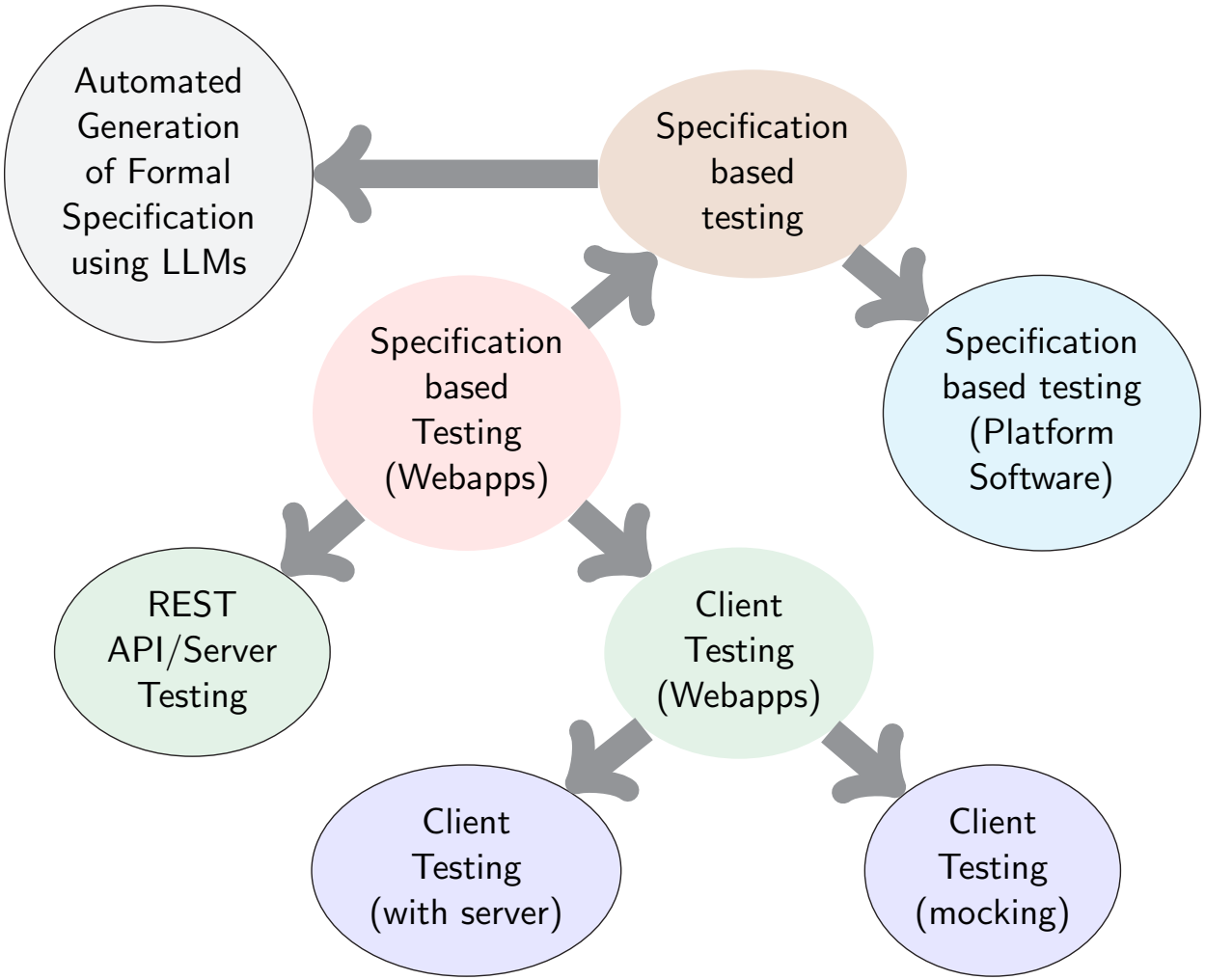
The TestGen Algorithm

A Roadmap



The TestGen Algorithm

A Roadmap



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Thank You!